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# Risk Analysis for Privacy in VANETs

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## We all know intuitively

- “Privacy is Important”
- “We need Changing Pseudonyms”

## But we don't know

- What kinds of attacks are probable
- What types of attackers will be there

## These are the results of a Risk Analysis

## Attack Trees

- Hierarchical
- Structured
- A common means to write down attacks [Schneier99]

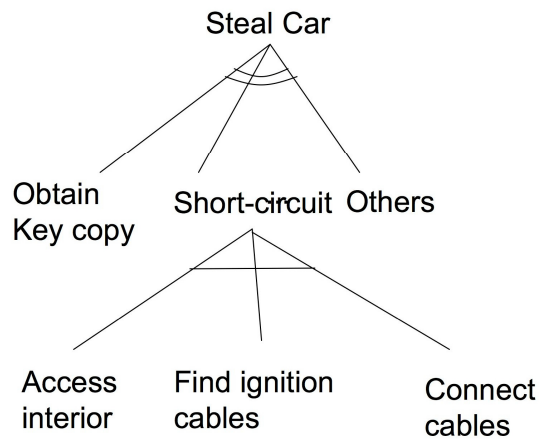
## Risk Analysis

- Mark up attacks as attack tree
- Estimate cost and other prerequisites for the attacker
- Find out cheap and probable attacks

## Assumptions

- Permanent pseudonym
- Broadcast of pseudonym, location, time

# Example Attack Tree



Goal **Steal Car** (OR)

- Obtain Key copy
- Short circuit (AND)
  - Access Interior
  - Find ignition cables
  - Connect cables
- Others

	graphical	textual
AND		Goal G0 (AND) G1 G2 ... Gn
OR		Goal G0 (OR) G1 G2 ... Gn

**AND** - all subattacks  
(cost add)  
**OR** - one subattack  
(cheapest cost counts)

## Violate privacy

- perception different for different people
- „obtain current or past location(s)“ according to Beresford's definition

## Two subattacks to violate privacy:

- Track location (pseudonym, time, location)
- Link pseudonym to name

## Finally: make use of obtained data

- matter of creativity

## Attack Method

- Build a grid of receivers
- Connect these receivers
- Store the data
- Process the data

## Coarse Cost Estimation (Examples)

- Road Side Unit (10 .. 30 EUR)
- Surveillance Camera (50 .. 100 EUR)
- DSL Connection (0 .. 15 EUR / Month)
- Storage (3K EUR / Tbyte)
- Access to database of phone provider (Proper Authorization)

## Attack Parameters

- Parameters for the abovementioned attack method
- Number of targets
- Coverage

## Attack Dimensions:

- All nodes, everywhere
- Some nodes, everywhere
- All nodes, some place
- Some nodes, some place

## Assumptions

- City of Berlin
- Permanent pseudonym
- Record a beacon every 3 seconds

## Conclusion

- Probably rather expensive for an attacker
  - Rather improbable attack
- Countermeasures:
    - Change the identifier
    - Do not provide more accurate data than necessary
    - Do not provide more data than necessary

## Build a grid of receivers

- Cheap, order of 200 K EUR

## Connect these receivers

- Expensive, order of 2 Million EUR per year

## Store the data

- 6 TB per day (this can probably be reduced)  
= 18 K EUR per day

## Process the data

- Not taken into account



# Some nodes, everywhere



## Possible attacks:

- All nodes, everywhere approach
- Use location requests, router functionality

## Attacker would go for second possibility

## Cost:

- A NOW Receiver (order of 100 EUR)
- Small database (Up to date PC)
- Dense enough network (Our Goal :- )

## I consider this as a probable attack.

## Countermeasures

- Artificially restrict max hops for location query
- Change pseudonym frequently
- Block frequent location queries by the network

Install a receiver at a fixed location

Log all beacons nodes at a specific location

Cost

- A NOW Receiver (order of 100 EUR)
- Small database (Up to date PC)
- Dense enough network (Our Goal :-)

I consider this as a probable attack.

Countermeasures

- Do you know any?

The last attack yielded only a pseudonym  
Useless without a name attached to it  
This attack is about getting your name  
Candidates:

- Use an existing database
- Restricted space identification
- Inference from external database information
- Ask

We proposed to create such a database to  
be able to revoke faulty/malicious nodes

This database links pseudonyms to names

Authorized entities may use this database

Cost:

- be authorized to use the database (high)

Not a very probable attack (except for the  
„authorized party“)

Use a publicly known one-to-one mapping of location and name (like home address)

Example:

- Obtain home location
- Obtain home address (trivial, using a map software)
- Get name from address (in theory, this data is contained in phonebooks)

Discussion:

- One-to-one mapping not always there
- For particular pseudonym, need full track to find this mapping

Some also call this *statistical disclosure*

## Method:

- Create (external) user profile
- Link the profile to the vehicle

## Example:

- Database says user drives blue car
- User just paid
- There is only one blue car on the parking (voilà)

## Discussion

- The better the profile or if there are additional tokens (such as parking-ticket, RFID tokens), this attack is easy

## Countermeasures

- Reduce accuracy of disclosed data
- Change pseudonyms shortly after/before possibly statistical disclosure

Use the name when embedded in a packet

Typical applications are

- Credit card payment
- Loyalty cards

Discussion

- The cheapest method to get the pseudonym name mapping
- Very probable
- But: user decides when to provide this (at least once: when installing the application)

The previous slides suggest that changing pseudonyms is a solution.

What are possible attacks on pseudonym change algorithms?

**Assumption:**

- Attacker has a set of messages with different pseudonyms (Track location classification applies)
- Objective: link messages coming from the same sender.

**Attack Classification**

- Based on non volatile data (e.g. vehicle brand)
- Protocol based attacks (e.g. beacon send period)
- Attacks based on physical parameters and constraints

**Cost Gain Analysis**

- Depends on the quality/quantity of available data
- Measured rather in terms of confidence in resolution (or meters after which a track is lost) than in money



Now the attacker knows your tracks (or parts thereof) and your name.

What could he do with it (Some ideas):

- Request a fine
- Blackmail
- Personalized advertisements, spamming
- Price discrimination
- Suspicion by location

Is current practice in commercial scenarios  
Change the price according to a profile of a  
user.

Example:

- Parking price is higher for cars of certain users
- Get the desired item (parking space, Big Mac, ...) only if you visited a certain location

Attacker:

- Commercial enterprises

Firefighter case (Based on a loyalty card profile a firefighter has been put in jail for six months)

Abuse of information

Information collected by commercial enterprises

Misused by authorities

Example

- anyone whose pseudonym has been observed at a crime scene may be guilty.

## This is still work in progress

### Attacks

- Global attacker is improbable but feasible, in particular if attackers team up ([www.payback.com](http://www.payback.com))

### Attackers

- Rather not governments (yet - access to the pseudonym - name database should be restricted),
- Rarely individuals
- Pretty surely commercial enterprises

### Countermeasures

- Changing pseudonyms are a good choice; make most attacks harder to carry out.

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